

Tutorial 6

P2 | Given: Alphabet A , blocklength $n \in \mathbb{N}$, $\mathcal{M} = A^n = \mathcal{C}$

a) An encryption function is an injective function $e_k : \mathcal{M} \rightarrow \mathcal{C}$, with $k \in \mathcal{K}$

injective function $f: f(x_1) = f(x_2) \Rightarrow x_1 = x_2$

$$|\{e_k(M) \mid M \in \mathcal{M}\}| = |\mathcal{M}|$$

$$|\{e_k(M) \mid M \in \mathcal{M}\}| \subseteq \mathcal{C} \Rightarrow e_k(\mathcal{M}) = \mathcal{C}$$

$\Rightarrow e_k$ is also surjective $\Rightarrow e_k$ is a bijection

A permutation π is a bijective function $\pi: X \rightarrow X$

$\Rightarrow \forall k \in \mathcal{K}$, the encryption e_k is a permutation with $X = A^n$

b) $A = \{0, 1\}$, $n = 6$, there are $N = 2^6 = 64$ elements

\Rightarrow there are $64! \approx 1.2689 \cdot 10^{89}$ different block ciphers.

P3/a) The bit error occurs in block $C_i, i > 0$, with block size BS

mode	M_i	max # errors	Remark
ECB	$E_K^{-1}(C_i)$	BS	Only block C_i is affected
CBC	$E_K^{-1}(C_i) \oplus C_{i-1}$	$BS+1$	M_i and one bit in M_{i+1} is affected
OFB	$C_i \oplus Z_i$	1	$Z_0 = C_0, Z_i = E_K(Z_{i-1})$
CFB	$C_i \oplus E_K(C_{i-1})$	$BS+1$	M_i and one bit in M_i
CTR	$C_i \oplus E_K(Z_i)$	1	$Z_0 = C_0, Z_i = Z_{i-1} + 1$

b) One bit is lost or added at block C_i at position $j, i > 0$

mode	block	position
ECB	i	1
CBC	i	1
OFB	i	j
CFB	i	j
CTR	i	j

In ECB and CBC all bits of all blocks $C_{i+t}, t \in \mathbb{N}_0$ may be corrupt.

In OFB, CFB, CTR, all bits beginning at position j in block i may be corrupt.

$$P1/k = (2D\ 61\ 72\ 69\ | \ 65\ 00\ 76\ 61\ | \ 6E\ 00\ 43\ 6C\ | \ 65\ 65\ 66\ 66)$$

$$= (w_0\ w_1\ w_2\ w_3)$$

$$a) k_0 = k = (w_0\ w_1\ w_2\ w_3)$$

$$b) k_1 = (w_4\ w_5\ w_6\ w_7)$$

Follow step 1 to calculate w_4 which are the first 4 Bytes of k_1

$$tmp \leftarrow w_3 = (65\ 65\ 66\ 66)$$

1) Evaluate Sub Bytes ($RotByte(tmp) \oplus Rcon(i/4)$)

$$1.1\ i/4 = 1; Rcon(1) = (Rc(1), 00, 00, 00) = (01, 00, 00, 00)$$

$$1.2\ RotByte(tmp) = (65\ 66\ 66\ 65) \quad // \text{cyclic left shift of bytes}$$

$$1.3\ SubBytes(65\ 66\ 66\ 65) = (4D\ 33\ 33\ 4D)$$

$$\Rightarrow tmp = (4C\ 33\ 33\ 4D) \quad // 1.3 + 1.1$$

$$2) w_4 \leftarrow w_0 \oplus tmp$$

	2	D	6	1	7	2	6	9
w_0	0010	1101	0110	0001	0111	0010	0110	1001
tmp	0100	1100	0011	0011	0011	0011	0100	1101
<hr/>								
Result	0110	0001	0101	0010	0100	0001	0010	0100
	6	1	5	2	4	1	2	4
tmp	4	C	3	3	3	3	4	D