

Polytechnic University

<http://wireless.poly.edu>



1/19

# Power and Data Rate Assignment for Max. Weighted Throughput in 3G CDMA: A Global Solution with Two Classes of Users

Virgilio Rodriguez, David J. Goodman, Zory Marantz

ECE Department  
Polytechnic University  
Brooklyn, NY  
email: [vr@ieee.org](mailto:vr@ieee.org)





# Problem description

- One “important” terminal shares a VSG-CDMA cell with several “ordinary” terminals.
- All terminals are delay-tolerant, and can operate at many data rates. They have “plenty of power”.
- Administrator wants to assign the data rate and transmit power of each terminal to maximize the cell **weighted** throughput.
- The throughput of the important terminal is weighted by  $\beta > 1$ .
- Weights admit various interpretations (price, utility, “priorities”, etc.)
- We knew that (i) **some** terminals should operate at maximal data rate, and (ii) **the others** should operate at a certain specific SIR value determined by the physical layer.
- We now determine how many terminals should be in each operating condition. This leads us to a global maximizer.





# Fundamental quantities of interest

- Terminals send data to a base station
- $R_c$ : chip rate ;  $R_i$  : data rate ;  $G_i = R_c/R_i$  : Spread Gain
- $G_i \geq G_0 \geq 1$  is required ( $R_i \leq R_{MAX} \leq R_C$ )
- $f_s(\gamma_i)$  : **probability of correct reception** of a **data packet**, in terms of received SIR. The properties of the physical layer (modulation, channel characteristics, FEC, diversity, etc.) are embodied into the FSF
- $\gamma_i := G_i \alpha_i$  is the SIR with  $\alpha_i$  the CIR given by

$$\alpha_i = \frac{h_i P_i}{\sum_{\substack{j=1 \\ j \neq i}}^N h_j P_j + \sigma^2} := \frac{Q_i}{\sum_{\substack{j=1 \\ j \neq i}}^N Q_j + \sigma^2}$$

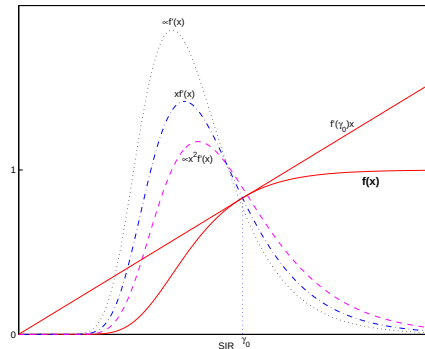
- $h_i$ : gain (“path loss”);  $h_i P_i := Q_i$ : received power
- Throughput of terminal  $i$ ,  $T_i(G_i, \alpha_i) \propto R_i f(G_i \alpha_i) \propto f(G_i \alpha_i)/G_i$





# An abstraction of the physical layer

- To accommodate most physical layers of interest, all that is assumed on the FSF is that it is a smooth “S-curve”:
  - “starts out” convex.
  - smoothly transitions to concave, as it approaches a horiz. asymptote.



- Ex: for non-coherent FSK with packet size  $M=80$ , no FEC, and perfect error detection, the FSF is  $f_s(x) = \left[1 - \frac{1}{2} \exp\left(-\frac{x}{2}\right)\right]^{80}$

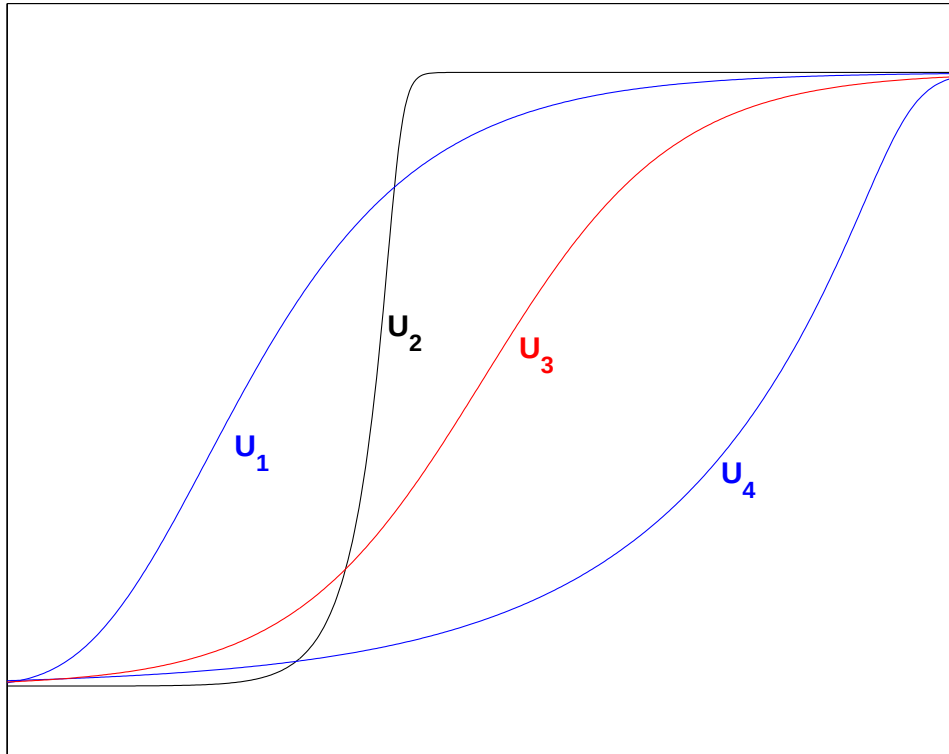


Figure 1: (1-FER) vs. SIR. S-curves can approximate “steps”, ramps, as well as concave and convex curves





# Optimization Model

- data rates can be determined through spreading gains,  $G_i$
- Power levels can be determined through power ratios. **BUT** CIR's need to be constrained so that they lead to feasible power levels

$$\max_{G_i, \alpha_i} \sum_{i=1}^{N-1} T_i(G_i, \alpha_i) + \beta T_N(G_N, \alpha_N)$$

subject to :

$$\sum_{i=1}^N \frac{\alpha_i}{1 + \alpha_i} = 1$$

$$G_i \geq G_0, i \in \{1, 2\}$$

- It is straightforward to write the Lagrangian for this problem and set up the associated first-order necessary optimizing conditions (FONOC)





# How many favored terminals?: Basic idea

- We know that (i) *some* terminals (“favored”) should operate at maximal data rate, and (ii) *the rest* should operate at the SIR  $\gamma_0$ , which maximizes  $f(x)/x$ .
- The SIR of the favored terminals, and the data rates of the non-favored must be chosen to satisfy FONOC.
- But how many in each group?
  - **First**, find, if available, a solution to FONOC with only the “important” terminal operating at maximal data rate (“**single favorite**” solution). **Then** seek a solution to FONOC in which both the important terminal **and** an ordinary one operate at maximal data rate. And so on.
  - After each solution is obtained, calculate the resulting throughput.
  - Choose the solution that gives the largest overall throughput





# Solving FONOC with one “favorite”

- To find the FONOC-solving SIR of the important terminal, find solutions to:

$$\phi\left(\frac{x}{G_0}\right) = \frac{1}{\beta}$$

- $\phi()$  is a “bell curve”, obtained from  $f$ . Therefore, this equation will have either 2 positive solutions, say  $x_1^*$  and  $x_2^*$ , or no solution at all (see fig).
- For any of these values that is greater than  $\gamma_0$ , obtain a corresponding  $\alpha$ , the FONOC-solving CIR for the non-favored terminals.
- The matching spreading gain is  $\gamma_0/\alpha$ . If  $\gamma_0/\alpha > G_0$ , a complete *feasible* solution to FONOC has been found, and the corresponding weighted throughput can be calculated.







# Solutions with many “favored” terminals

Let  $x$  and  $y$  be the SIR of, respect., the important and ordinary terminals operating at maximal data rate.

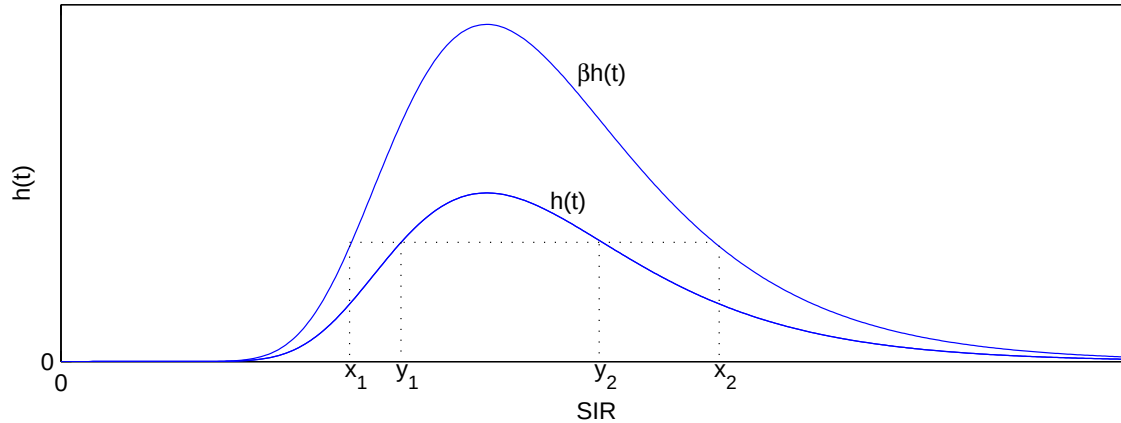
- To obtain  $x$  and  $y$  we must solve a system of 2 non-linear equations of the form:

$$\beta h(x) = h(y) \text{ and } \pi(x, y) = 0$$

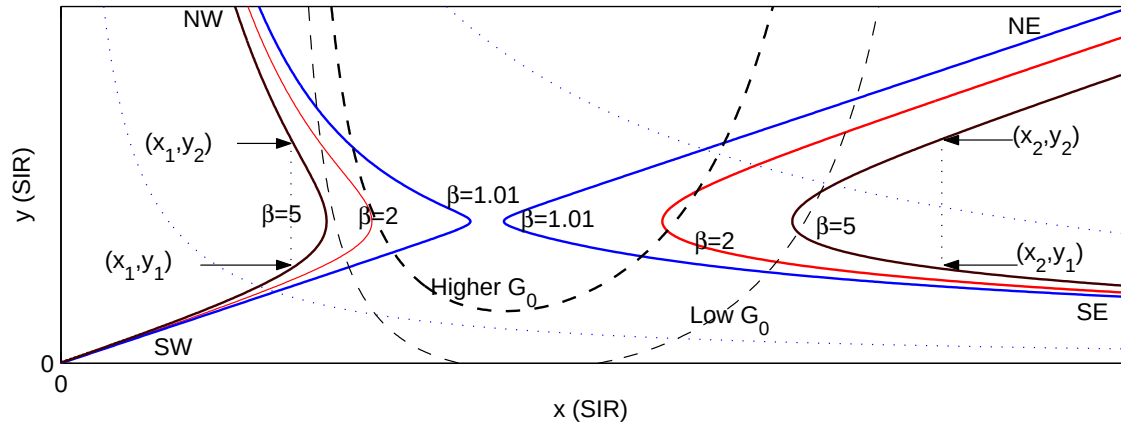
- $h()$  is a “bell curve” obtained from  $f$ . The pairs  $(x, y)$  that satisfy  $\beta h(x) = h(y)$  give rise to an “X-shaped” graph.
- $\pi(x, y) = 0$  follows directly from the constraint on the CIR’s ( $\Sigma() = 1$ ). and give rise to a U-curve, which turns into an L-curve (hyperbola) when all terminals operate at maximal data rate.
- The 4 intersections between the X and U (or 2 between the X and the L) are candidates for (local) maximizers. Some of them can be eliminated by applying some additional conditions.



$(x_1, y_1)$ ,  $(x_2, y_2)$ ,  $(x_1, y_2)$  and  $(x_2, y_1)$  are possible solutions to  $\beta h(x)=h(y)$



"X-shaped" graphs showing points  $(x,y)$  satisfying  $\beta h(x)=h(y)$



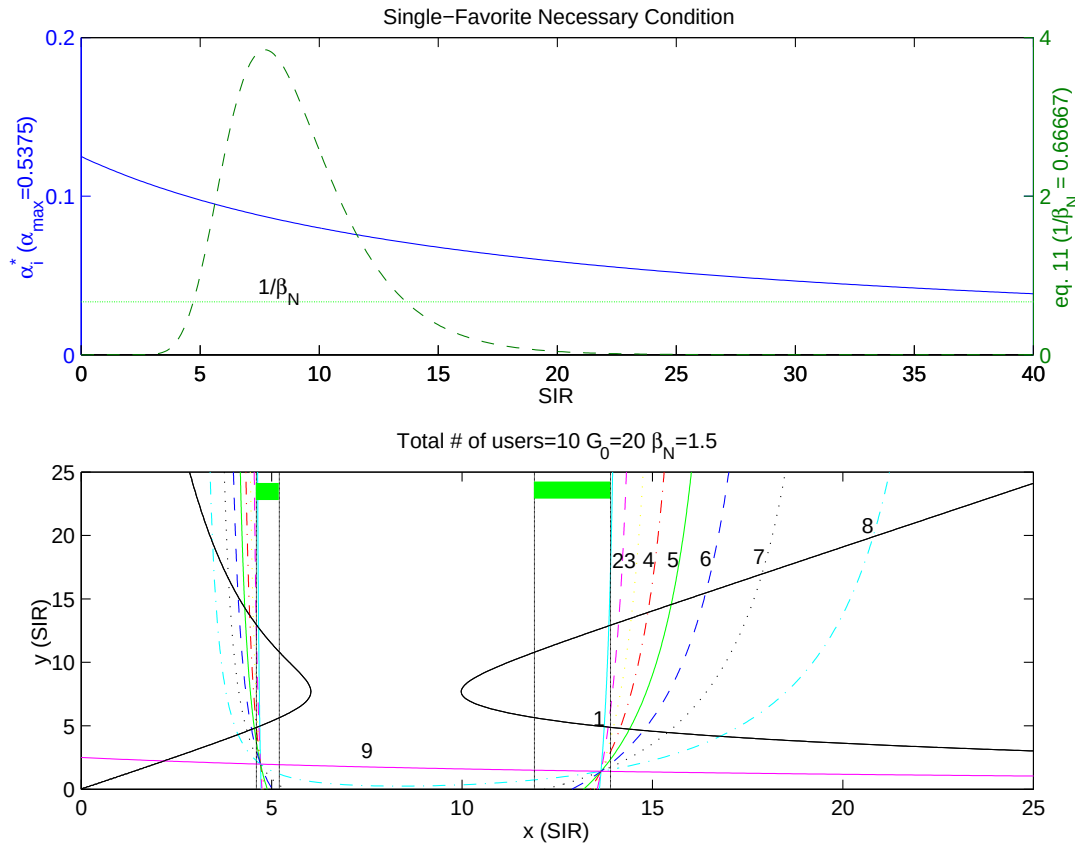


Figure 2: Single-favorite solution, at  $x \approx 13$  (top), yields a *weighted* throughput of  $0.12R_c$ . DFBS at  $(13.8, 12.8)$  maximizes the weighted throughput at  $0.7R_c$ . The “all favored” solution ( $n_1 = 9$ ) leads to a minimum.

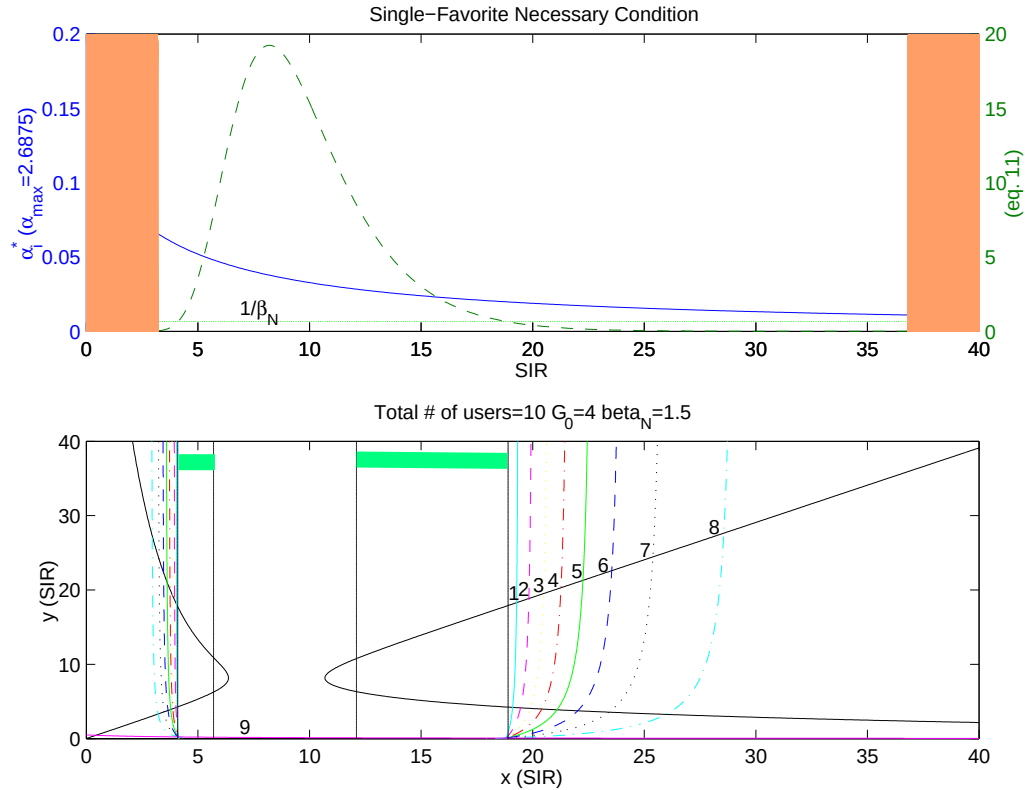


Figure 3: With a small  $G_0 = 4$  and a moderate  $\beta = 1.5$  the SFBS exists (top), and leads to the maximum. But all the multi-favorite solutions fail. The “all favored” solution leads to a minimum.



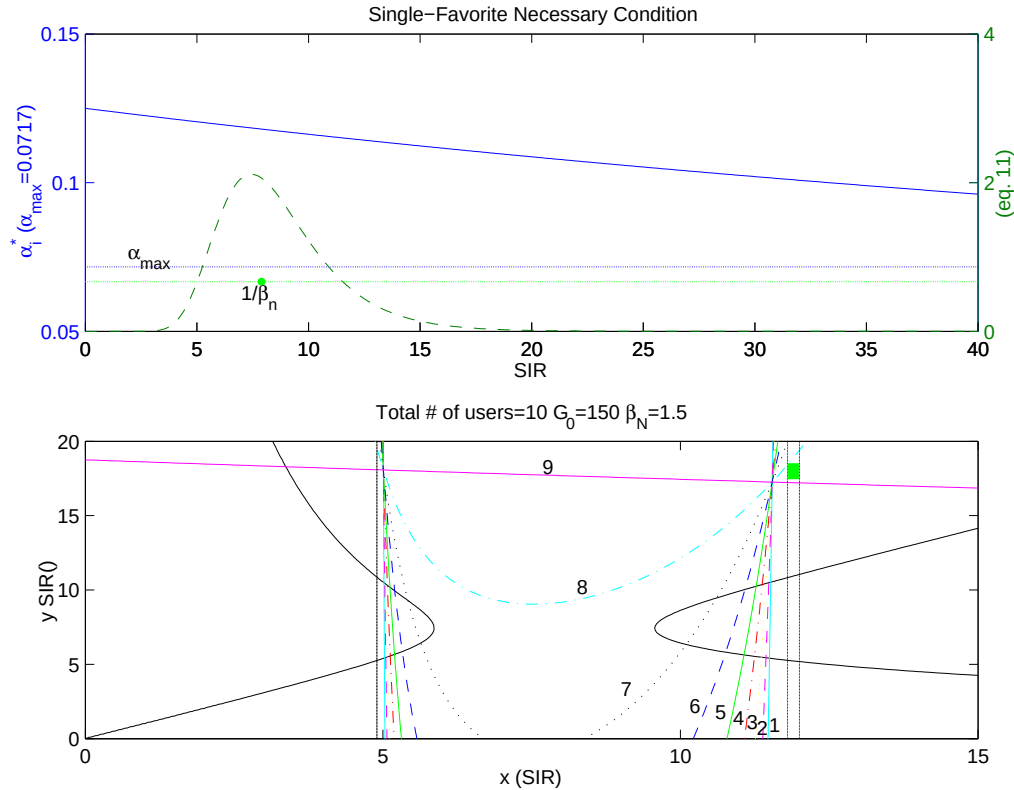


Figure 4: With a high  $G_0 = 150$  and  $\beta = 1.5$  no SFBS is available because for any  $x$ ,  $\alpha$  will be too high (its matching spreading gain will fall below  $G_0$ ). Multi-favorite solutions with  $1 \leq n_1 \leq 8$  also fall outside the acceptable range. The “all-favored” solutions leads to the global maximizer.





# Discussion

- We have allocated data rate and power for maximal **weighted** throughput in a 3G CDMA (VSG) cell. Weights may reflect “importance”, utility, or price paid per bit delivered.
- To accommodate most physical layers, we assume that **all that is known** about the frame-success function (FSF) is that it is an “**S-curve**”.
- Each physical layer has a “preferred” SIR, at the “knee” of the FSF.
- (i) **Some** terminals (“favored”), including the important one, should operate at maximal data rate, and (ii) **the others** should operate at the SIR “preferred” by the physical layer.
- With several terminals operating at the highest data rate, the intersection of an X-shaped graph and a U or L-shaped graph (from SIR-feasibility constraint) determines the SIR of the “favored” terminals.
- We have presented a procedure to determine how many terminals should operate at maximal data rate; this has led us to the global maximizer.





# Extensions and future directions

- We have already started to extend this analysis to include noise (out of cell interference), and the presence of power-limited media terminals (CISS-04).
- This analysis can be further extended to consider
  - QoS requirements
  - power limits
  - many cells





# Technological Motivation : VSG-CDMA

- Modern (3G) wireless nets are expected to accommodate terminals operating at very different data transmission rates.
- Variable Spreading Gain CDMA can accommodate multi-rate traffic, and is supported by current 3G standards.
- In a VSG CDMA system, chip rate is common, but each terminal's spreading (processing) gain is the ratio of the common chip rate to the terminal's bit rate.







# Outline of research program

- Analytical Core
- Single-user applications: choose power, data rate and/or coding rate for data, image, video transmission
- Decentralized multi-user applications:
  - Game formulation
  - Mechanism design
- Centralized data throughput maximization
  - Without noise
  - With noise and media terminals present





# Overview of our analytical framework

- Many radio-resource optimizations share a common **analytical core**, which enables **robust** and **tractable** analysis and provides **clear answers** in fairly **general** scenarios
- It involves
  - A tractable abstraction of the physical layer
  - A tractable abstraction of the human visual system
  - A fundamental result: **maximize  $f(x)/x$  with  $f$  an “S-curve”**.
- Problems to which this framework applies:
  - Power and coding rate choice for media files (images, video)
  - Choosing the “right amount” of media distortion
  - Decentralized power control for 3G CDMA
  - Data rate and power allocation for maximal cell throughput when data and media terminals share a CDMA cell





THANK YOU!!!

WIRELESS.POLY.EDU

