Simple decentralised market-oriented allocation of sub-channels and power for access-point to terminal multi-carrier communication

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Market-based OFDMA resource management 1/

## **Executive Overview**

- Orthogonal frequency-division multiple-access (OFDMA) is the technology of fourth generation (4G) cellular networks
- Efficient resource allocation is critical to 4G success
- Main issues are allocating sub-channels and power levels to individual terminals
- A decentralised solution is preferable because of the high complexity of finding the globally optimal allocation
- For the base-station-to-mobile link, we combine auctions
  to allocate the sub-channels with per-Watt pricing
- Our scheme can involve real money (service fees), or be implemented as a decentralised low-complexity algorithm
- We now discuss conceptual/qualitative/analytical issues. Numerical studies focused on 4G networks are anticipated

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## The free market as a paradigm for algorithms

- A complex technological system can be "efficiently" managed as an "economy"
- The system can be viewed as integrated by many "agents"
- Agents may correspond to actual human beings, or may be machines, or processes
- The system administrator sets up some relatively simple rules for resource use and behaviour (prices, auctions, rewards, punishments, etc)
- Each agent behaves and utilises resources as an economic entity seeking to maximise its "preferences" while obeying the rules and budget constraints (energy, power, bandwidth, etc).
- If the rules are "right", the complex system produces "efficient" results

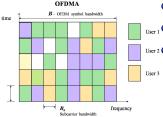
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#### Auctions for short-term allocations: Why?

- To quickly allocate resources to those that most value them
- A terminal's valuation of a resource could either
  - represent the "true" "willingness to pay" of a user, or
  - be a "priority" index computed/adjusted by software
- The Dutch auction format is of interest:
  - Public "clock" displays a progressively falling price
  - Eventually bidder that most values object "takes it"



#### **OFDMA** management: Basic issues



- OFDMA downlink cell with K terminals
- Bandwidth split into N sub-carriers
- User 1 Resources assigned once per time-slot
- User2 Channel-state info at BS enables:
  - Dynamic sub-carrier assignment (NP-hard!!)
  - Adaptive power allocation ("water filling")

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- Adaptive modulation and coding
- Allocate "sub-channels" (several sub-carriers) to reduce overhead

## Economorphic OFDMA management

- Semi-decentralised market oriented (economic) approach
- Basic idea:
  - Just before the time-slot of interest, network auctions the sub-carriers in parallel by means of a descending-price "clock" auction (Dutch auction)
  - When the price for a given sub-carrier is low enough, the terminal that "most values" it sends a buying signal.
  - This terminal also purchases its optimal amount of down-link power for that sub-carrier, given a per-Watt price set by the network
  - After a given sub-carrier is won, the price of each of the others continue to drop until a terminal buys it.

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## Initial bid calculations

- A terminal's bid closely related to its "valuation" of a SC
- Valuation: Benefit minus associated cost
- "Benefit": "value" of the (additional) information that the terminal can transfer if it wins that SC
- Assuming one sub-carrier per sub-channel, a data (delay tolerant) terminal that has NOT yet won anything, values SC *n* as  $\beta_i R_{i,n}(p_{i,n}^*) c_i p_{i,n}^*$  where :
  - *β<sub>i</sub>* is the monetary value of one information bit successfully transferred multiplied by the SC bandwidth
  - $R_{i,n}(p) = \log_2(1 + ph_{i,n})$  (information bits/Hertz terminal can transfer over SC *n* with power *p* and channel-gain over noise equal to  $h_{i,n}$ )
  - For price  $c_i$ ,  $p_{i,n}$  is the solution to single-variable equation  $R'_{i,n}(p) = c_i/\beta_i$  and  $p^*_{i,n} = \min(p_{i,n}, P)$  where *P* is the remaining power.

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## Bid re-calculation after a previous winning

 If terminal *i* wants to evaluate SC *n* after it has won SC *m* it must solve:

$$\max_{x,y} \beta_i(R_{i,m}(x) + R_{i,n}(y)) - c_i(x+y) \quad \text{s.t. } x+y \le P$$

- If  $x^*, y^*$  are the optimisers, the terminal only needs  $x^* + y^* p_{i,m}^*$  because it bought  $p_{i,m}^*$  when it won SC *m*.
- The increase in benefit brought by sub-channel n is:

$$\beta_i(R_{i,m}(x^*) + R_{i,n}(y^*) - R_{i,m}(p_{i,m}^*))$$

( the benefit from having both m and n minus that of having m alone).

• Thus, the valuation of SC *n* after *m* has been won is:

$$\beta_i(R_{i,m}(x) + R_{i,n}(y) - R_{i,m}(p_{i,m}^*)) - c_i(x + y - p_{i,m}^*) \equiv [\beta_i(R_{i,m}(x) + R_{i,n}(y)) - c_i(x + y)] - [R_{i,m}(p_{i,m}^*)) - c_i p_{i,m}^*]$$

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## Valuation after several winnings

● A terminal that has won *M* − 1 bands and evaluates an additional one must solve a problem of the form:

$$\max_{x_{1},\dots,x_{M}} \beta_{i} \sum_{m=1}^{M} R_{i,m}(p_{m}) - c_{i} \sum_{m=1}^{M} p_{m}$$
  
s.t. 
$$\sum_{m=1}^{M} p_{m} \leq P$$
$$p_{m} \geq 0$$
(1)

 The solution (provided separately), has the "water-filling" form: p<sub>m</sub>+1/h<sub>m</sub> = 1/(c+λ)

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#### Implementation issues

- Auction requires tight synchronisation among terminals, i.e., a "common clock" (NOT a problem for 4G networks)
- Auction parameters (initial price, the clock "tick", and price "step") should be chosen judiciously, considering statistics of the terminals' "valuations", and signal processing/travel times, among other factors
- Simultaneous bids are in principle possible. However, if channel gains are continuous random variables, so are valuations and bids; therefore, the probability of tied bids is negligible.
- Can this scheme be implemented within the constraints imposed by present 4G standards?

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#### Summary

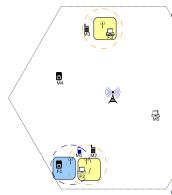
- Decentralised sub-channel/power allocation scheme for the forward-link of an OFDMA cell.
- Based on simple descending-price auctions, one for each sub-channel, run in parallel at each time slot
- Scheme is low complexity: global NP-hard sub-channel/power allocation problem avoided.
- Key computation done by each terminal in calculating its bid in a "channel by channel" basis.
- Bid calculation become more complex, as a given terminal wins more sub-channels
- Tendency: sub-channel goes to terminal that can transfer the "highest value" of information through it
- Auction can involve real money (service fees), or the bids may represent "priority" indeces
- Data traffic ("best effort") (e-mail, web browsing, etc) has been assumed. If a terminal requires minimum data rate, it must bid "high enough" to always win necessary resources
- So far work is conceptual/analytic. Simulation anticipated.

# Extension: OFDMA Femtocells



- femto access point (FAP): low-cost, low-power device suitable for home-owner installation
- operates underneath a standard cell, in licensed spectrum bands and can
  - improve user experience in indoor locations,
  - increase overall system capacity
  - with modest monetary investment.
- But the unplanned dynamic nature of these cells, significantly complicates resource allocation and interference control.

# Extension: simple management of OFDMA Femtocells



- Our scheme can be adapted to femtocells
  - Interference controlled through confirmation messages, e.g.,
    - if M1 or M2 hear FAP-2 confirm a buy message, M1 and M2 infers the SCI is no longer available
    - BS confirms over the air confirmation, leaving available BS power unchanged .
    - Any FC terminal that did not hear FAP-2
      e.g., FAP-3 will continue to treat the SC as available
- See: IWCMC starting June'10, in France